

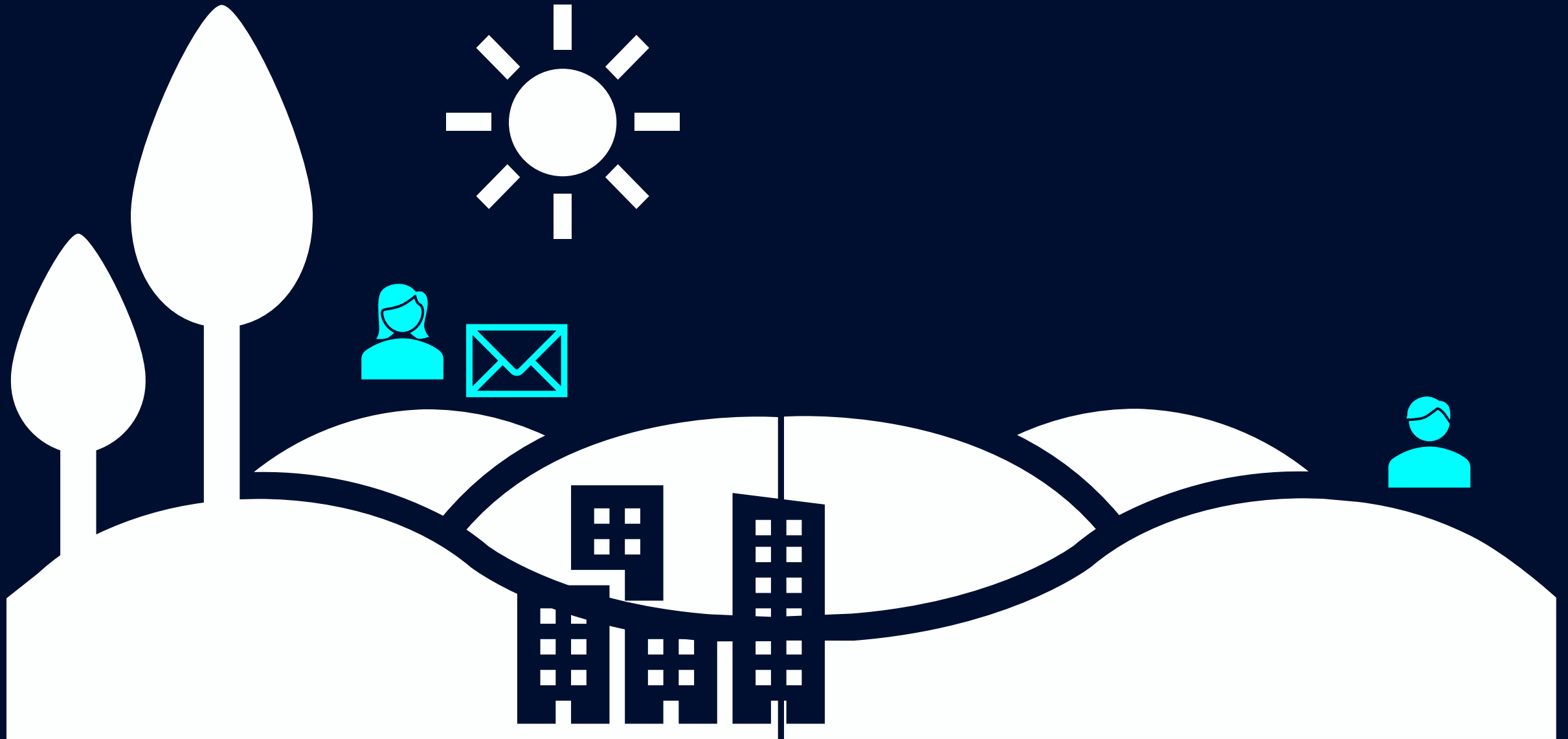
Cracking the codes of Alice

Tjitske Koster | PhD at TU Delft

2-3-2026



Alice and Bob



Eve



How to hide
the message?



Tjitske Koster

TU Delft






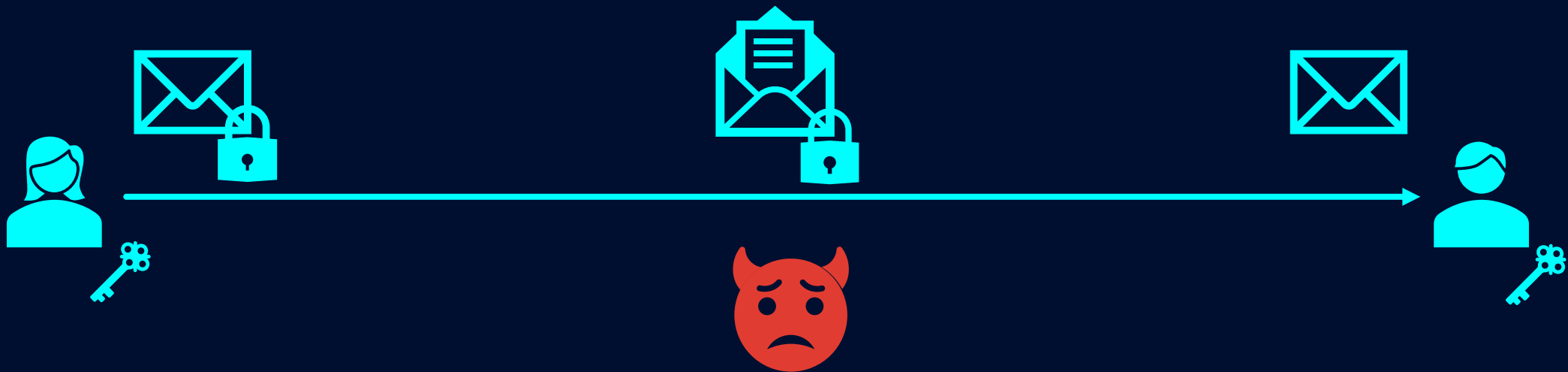
Goal today

- Introduce the endless battle of cryptography

Goal of Alice

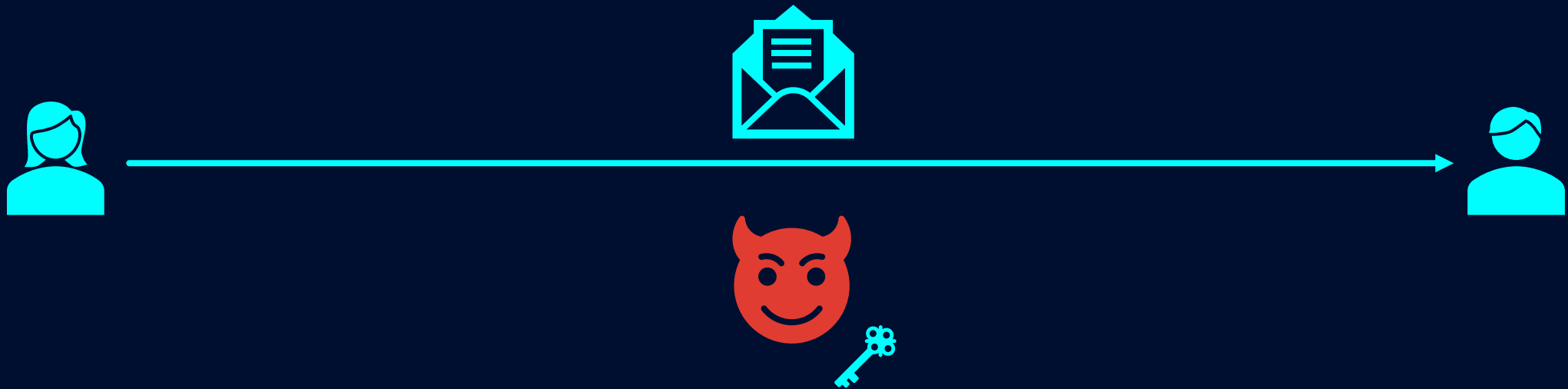


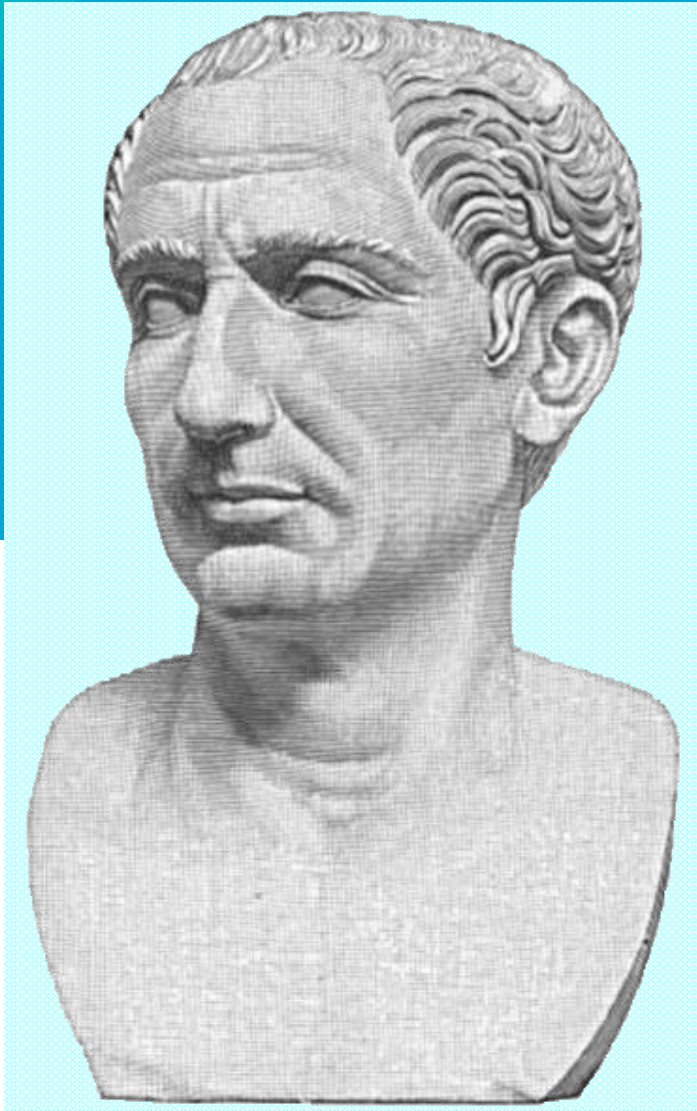
- A code  that only Alice and Bob know
- Alice can *Encrypt* the message  **Ciphertext**
- Bob can *Decrypt* the message  **Plaintext**
- If Eve sees the message, she will not understand



Goal of Eve (us)

- Break the code of Alice





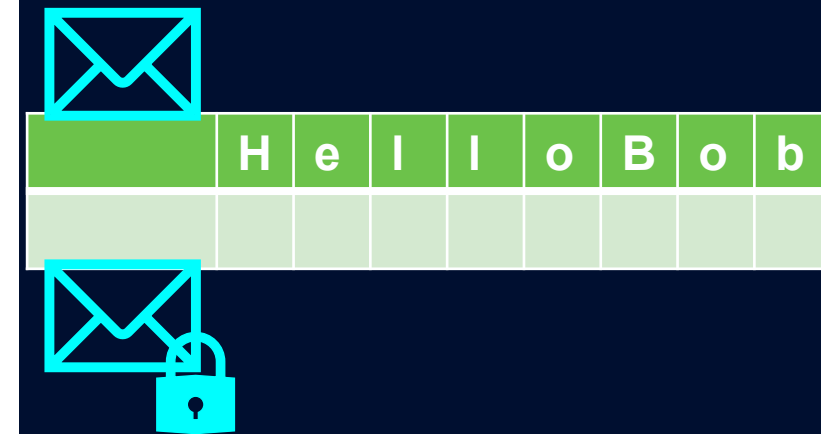
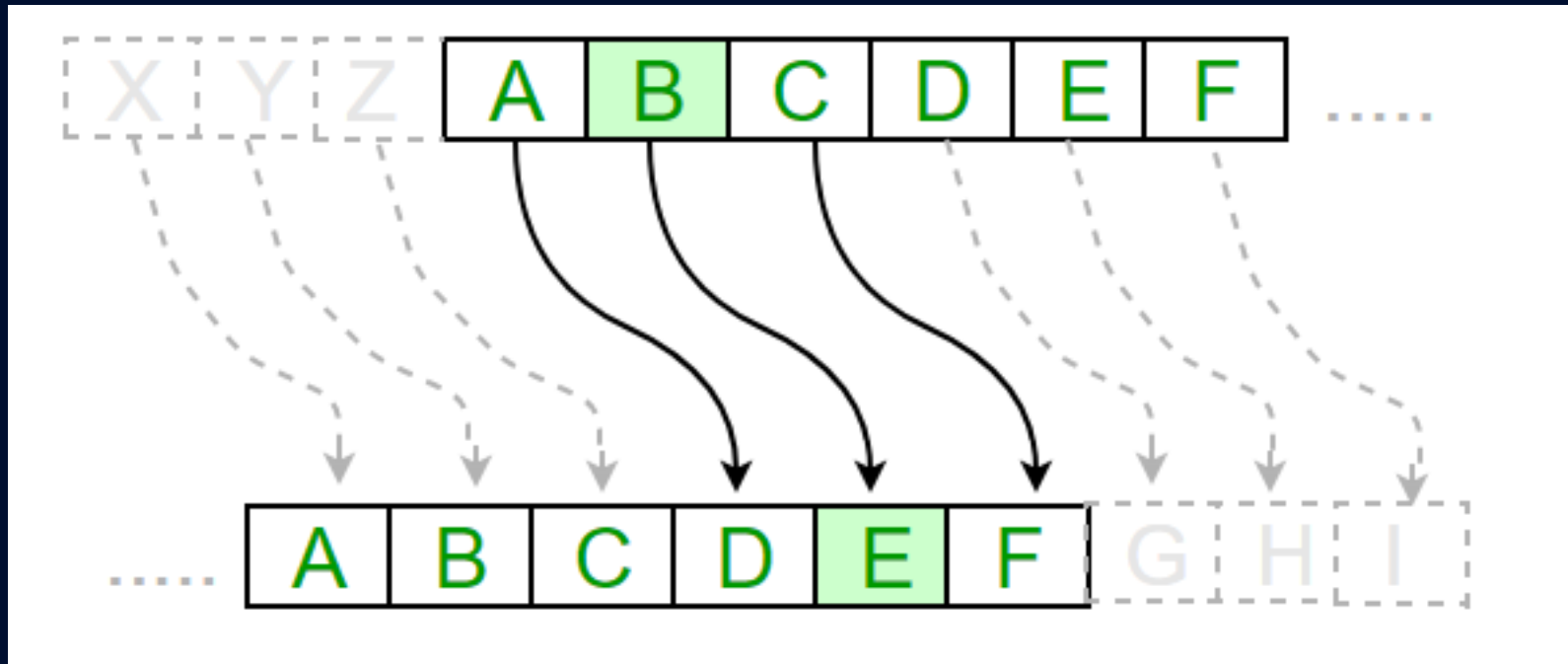
Caesar cipher

50 BC

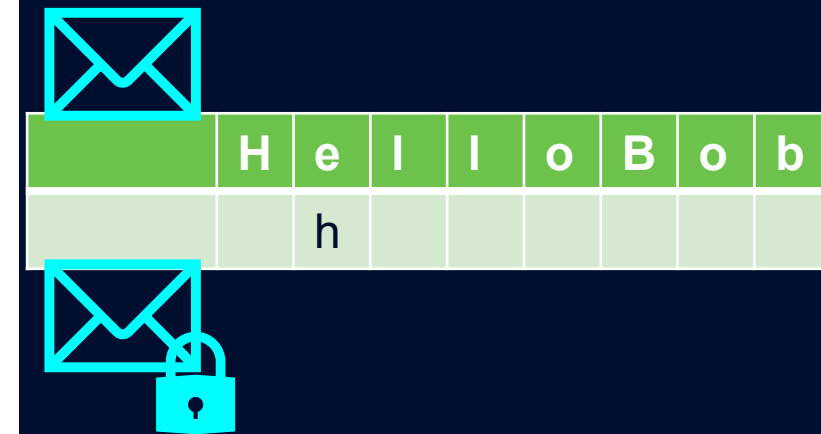
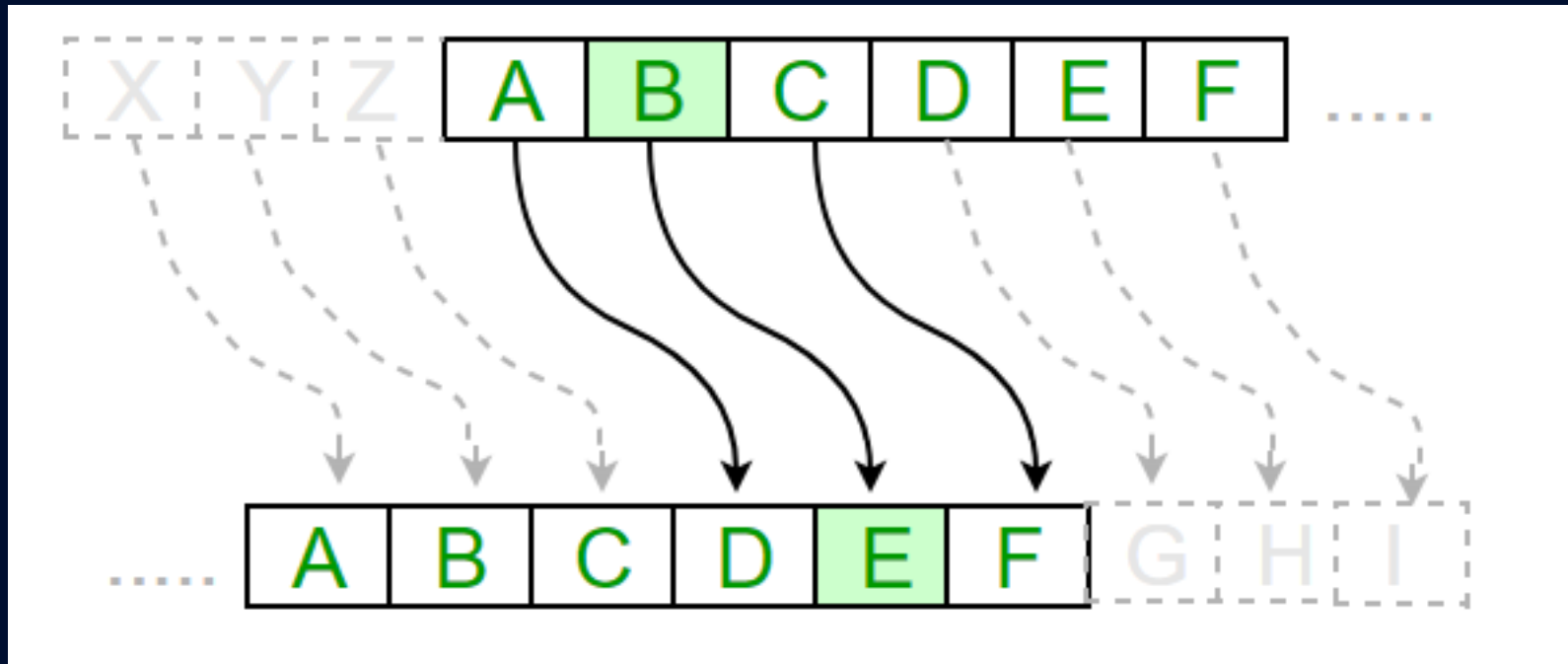
The first encryption Caesar Cipher



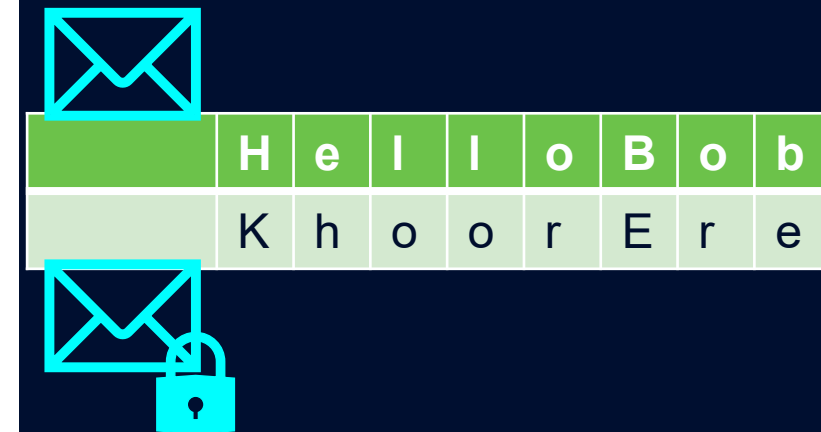
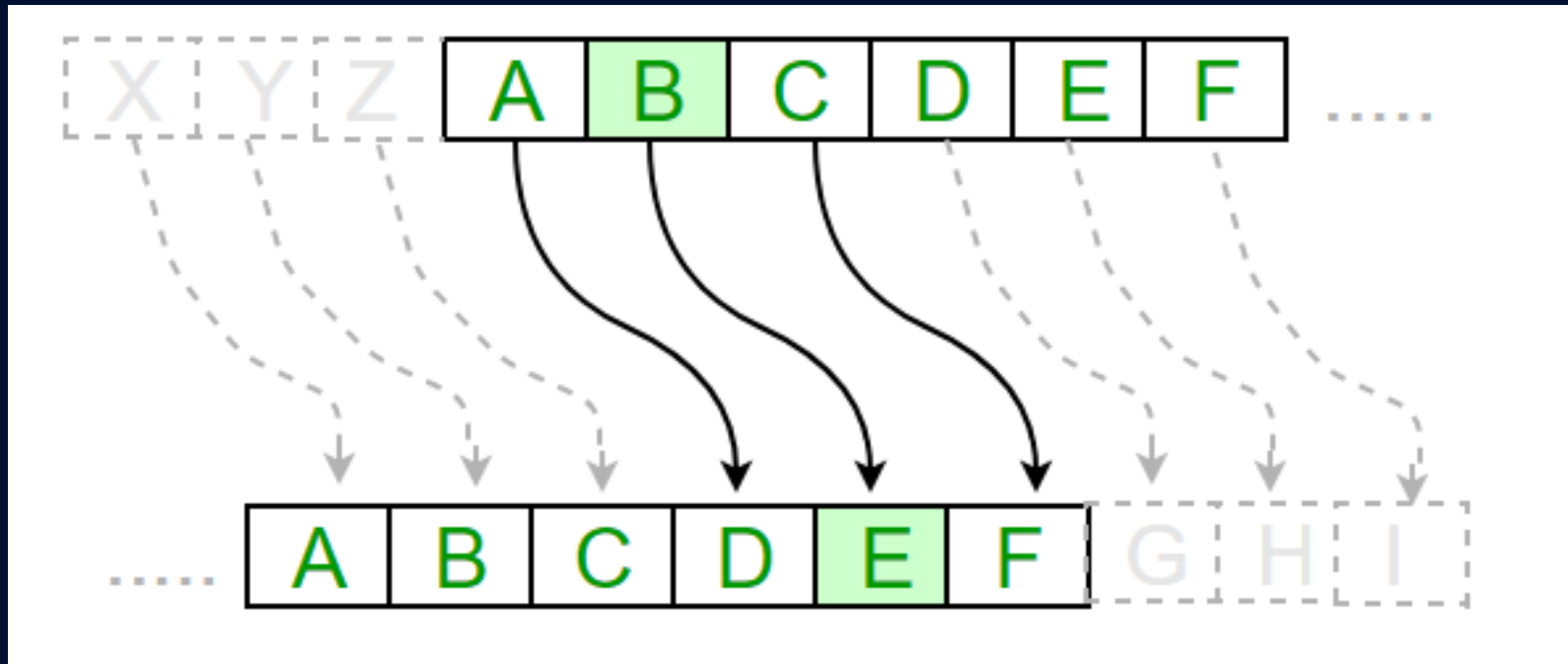
- Replace each letter with a letter K places further in the alphabet



The first encryption Caesar Cipher



The first encryption Caesar Cipher



Breaking Caesar

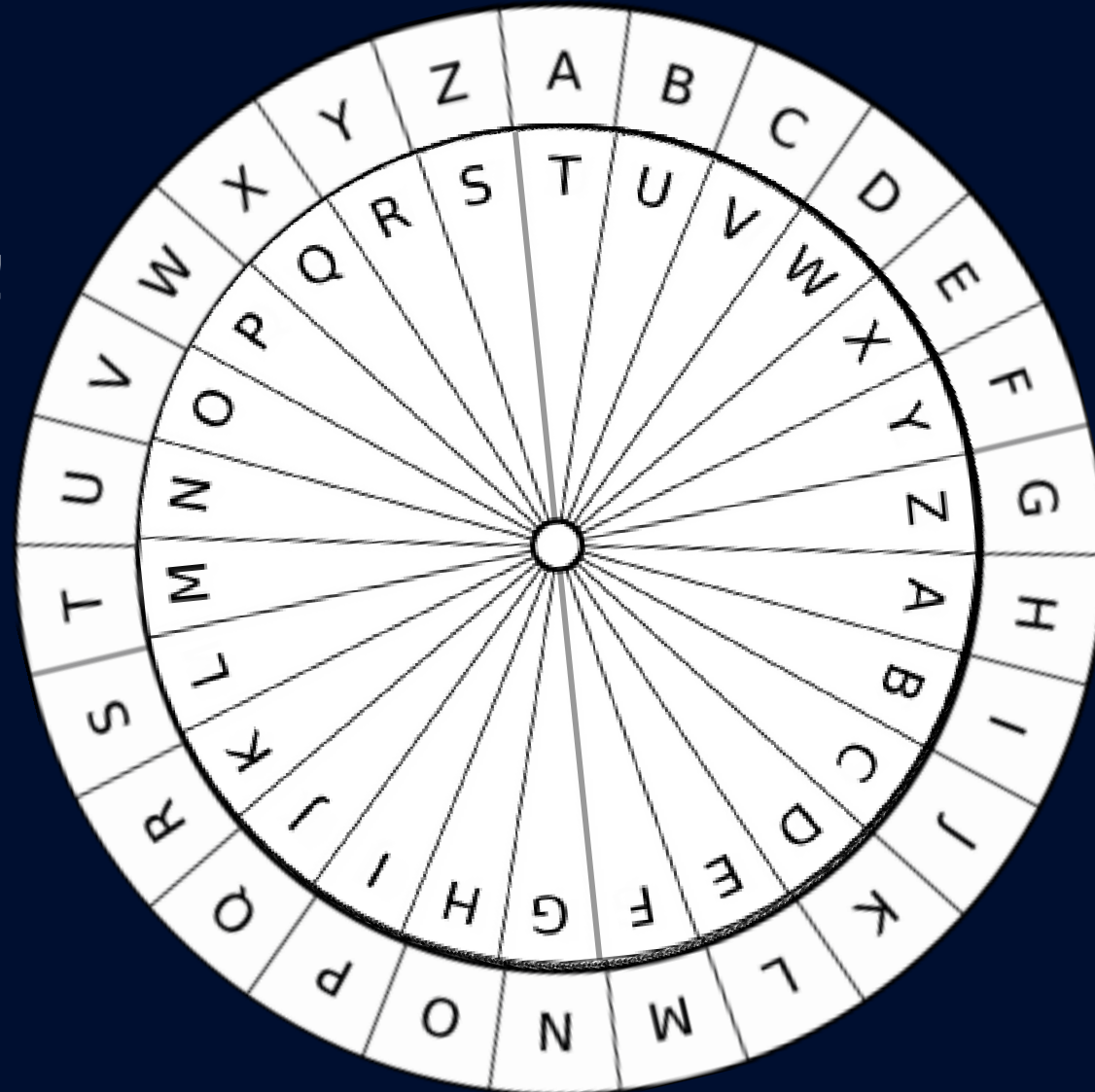


Olssv ivi!

Breaking Caesar



Olssv ivi!



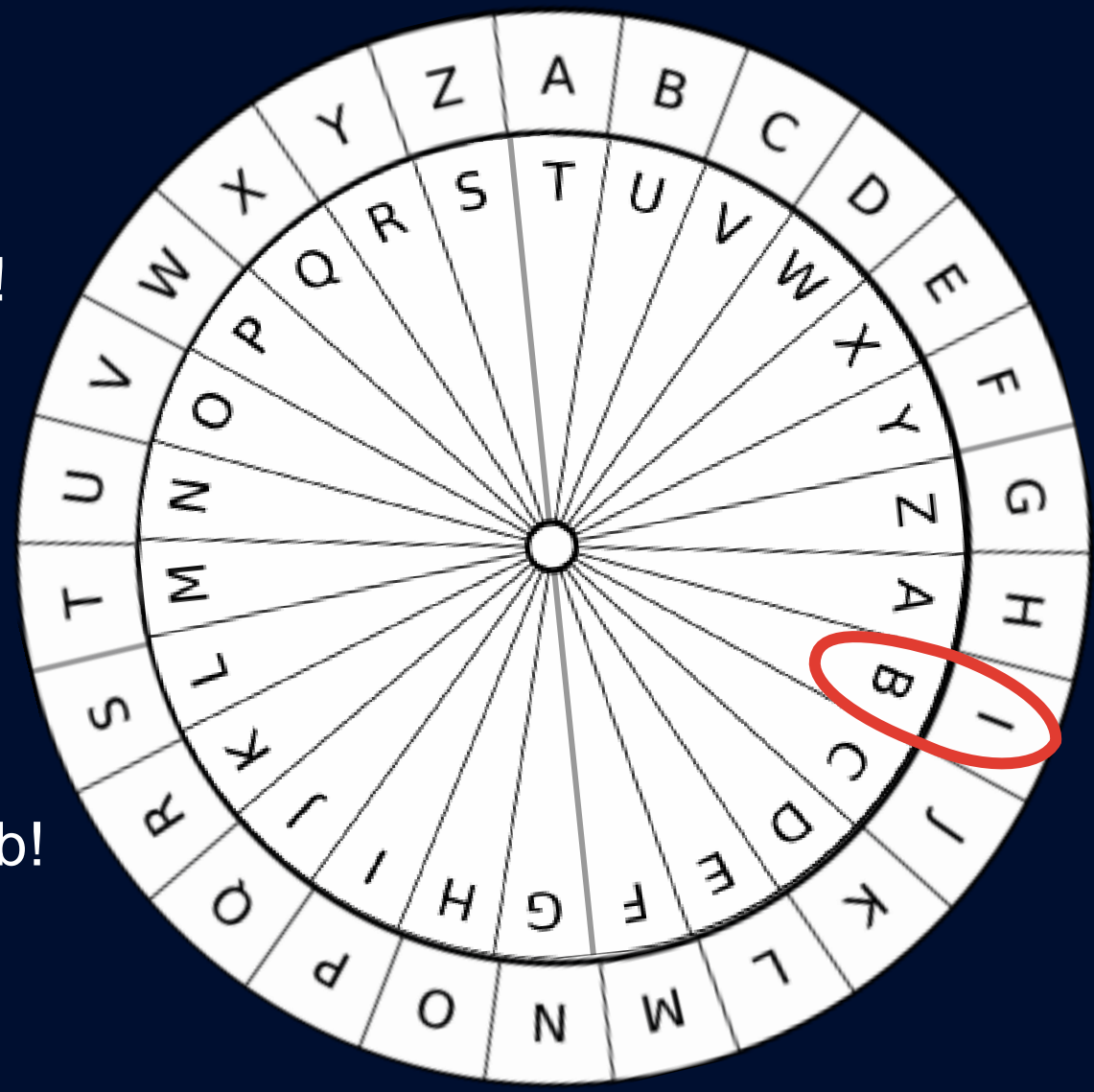
Breaking Caesar



Olssv ivi!



Hello Bob!



Linguistic analysis



Uvd kbyapun aopz aptl Zohoyhghk ohk ivyul rpun
Zohoypfhy aoyll zvuz. Vu aol aovbzhuk huk mpyza
upnoa, dolu zol ohk luklk aol ahsi vm Th'hybm, zol yvzi
huk rpzzlk aol nyvbuk ilmvyi opt, zhfpun: “nylha rpun,
mvy h aovbzhuk huk vul upnoaz P ohcl illu yljvbuapun
av fvb aol mhisiz vm whza hnlz huk aol slnlukz vm
hujplua rpunz. Thf P thrl zv ivsk hz av jyhcl h mhcvby
vm fvby thqlzaf?”
Lwpsvnl, Ahsiz myvt aol Aovbzhuk huk Vul Upnoaz



Linguistic analysis

Uvd kbyapun aopz aptl Zohoyhghk ohk ivyul rpun
Zohoypfhy aoyll zvuz. Vu **aol** aovbzhuk huk mpyza
upnoa, dolu zol ohk luklk **aol** ahsi vm Th'hybm, zol yvzi
huk rpzzlk **aol** nyvbuk ilmvyi opt, zhfpun: “nylha rpun,
mvy h aovbzhuk huk vul upnoaz P ohcl illu yljvbuapun
av fvb **aol** mhisiz vm whza hnlz huk **aol** slnlukz vm
hujplua rpunz. Thf P thrl zv ivsk hz av jyhcl h mhcvby
vm fvby thqlzaf?”
Lwpsvnl, Ahsiz myvt **aol** Aovbzhuk huk Vul Upnoaz



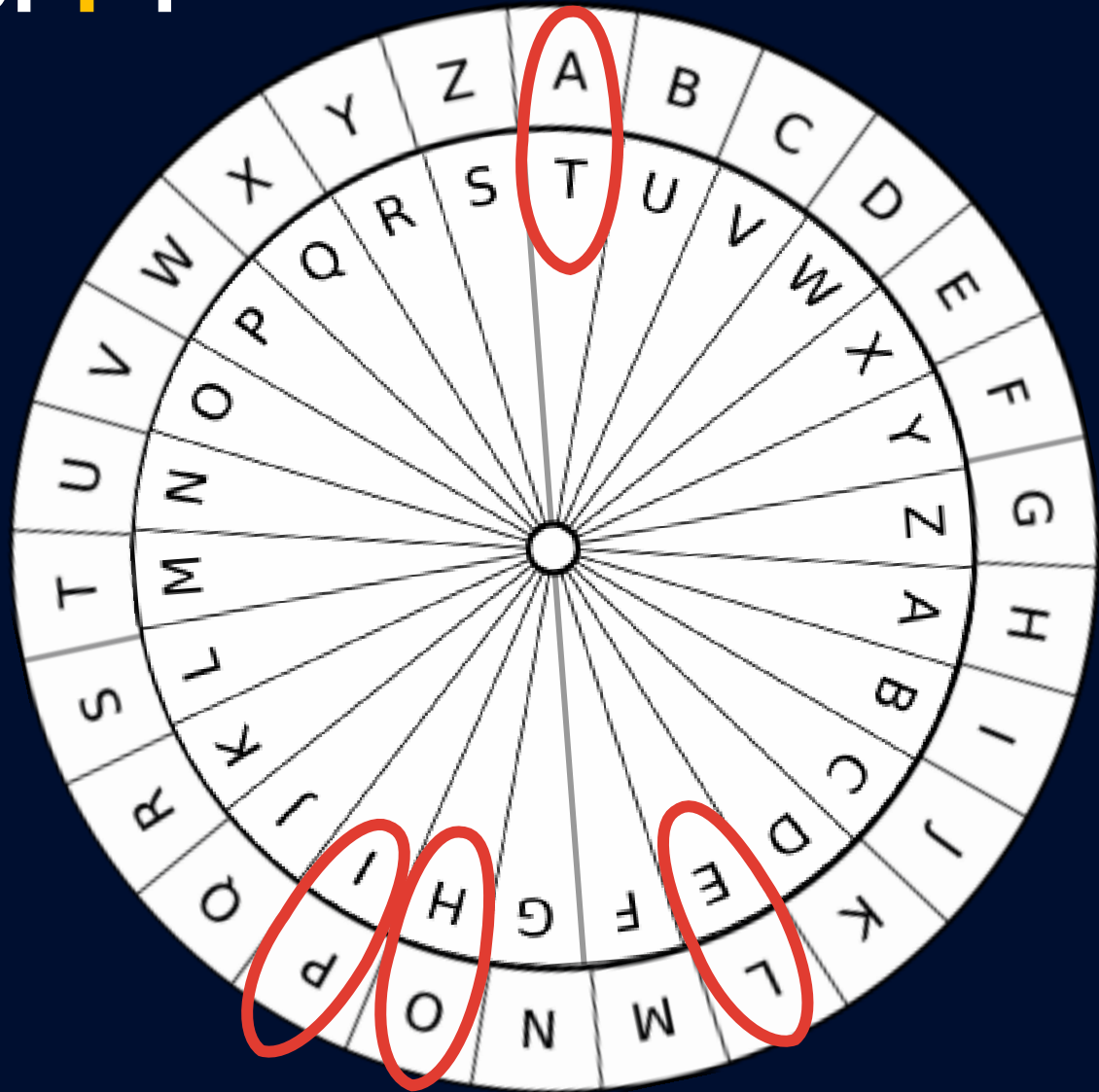
Linguistic analysis

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mvy h aovbzhuk huk vul upnoaz **P** ohcl illu yljvbuapun
av fvb **aol** mhisiz vm whza hnlz huk **aol** slnlukz vm
hujplua rpunz. Thf **P** thrl zv ivsk hz av jyhcl h mhcvby
vm fvby thqlzaf?”
Lwpsvnl, Ahsiz myvt **aol** Aovbzhuk huk Vul Upnoaz

What would **aol** be, or **P**?

- aol → the ?
- P → I?

- A → T
- O → H
- L → E
- P → I

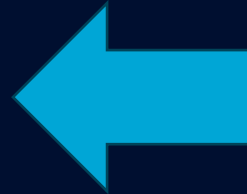


Linguistic analysis - Decryption



Now during this time Shahrazad had borne king Shahriyar three sons. On the thousand and first night, when she had ended the tale of Ma'aruf, she rose and kissed the ground before him, saying: “great king, for a thousand and one nights I have been recounting to you the fables of past ages and the legends of ancient kings. May I make so bold as to crave a favour of your majesty?”

Epilogue, Tales from the Thousand and One Nights



Uvd kbyapun aopz aptl Zohoyhghk ohk ivyul rpun Zohoypfhy aoyll zvuz. Vu **aol** aovbzhuk huk mpyza upnoa, dolu zol ohk luklk **aol** ahsI vm Th'hybm, zol yvzl huk rpzzlk **aol** nyvbuk ilmvyI opt, zhfpun: “nyIha rpun, mvy h aovbzhuk huk vul upnoaz **P** ohcl illu yIjvbuapun av fvb **aol** mhisIz vm whza hnlz huk **aol** slnlukz vm hujplua rpunz. Thf **P** thrl zv ivsk hz av jyhcl h mhcvby vm fvby thqlzaf?”
Lwpsvnbl, AhsIz myvt **aol** Aovbzhuk huk Vul Upnoaz

What went wrong?



- Spaces
- Capitals
- Punctuation

- Only 25 encryptions possible; **brute force attack** is feasible



Vigenère cipher

16th century

The second encryption



Vigenère



- We assign numbers to each letter

A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9

K	L	M	N	O	P	Q	R	S	T
10	11	12	13	14	15	16	17	18	19

U	V	W	X	Y	Z
20	21	22	23	24	25

The second encryption



Vigenère



- Alice and bob agree on a secret: **KEY**
- Alice translate message to numbers
- Alice writes the secret
- Alice translate secret to numbers
- Alice adds the message to the secret
- Alice translate numbers to letters



H	E	L	L	O	B	O	B
07	04	11	11	14	01	14	01
K	E	Y	K	E	Y	K	E
10	04	24	10	04	24	10	04
17	8	9	21	18	25	24	05
R	I	J	V	S	Z	Y	F



$$7 + 10 = 17$$

$$4 + 4 = 8$$

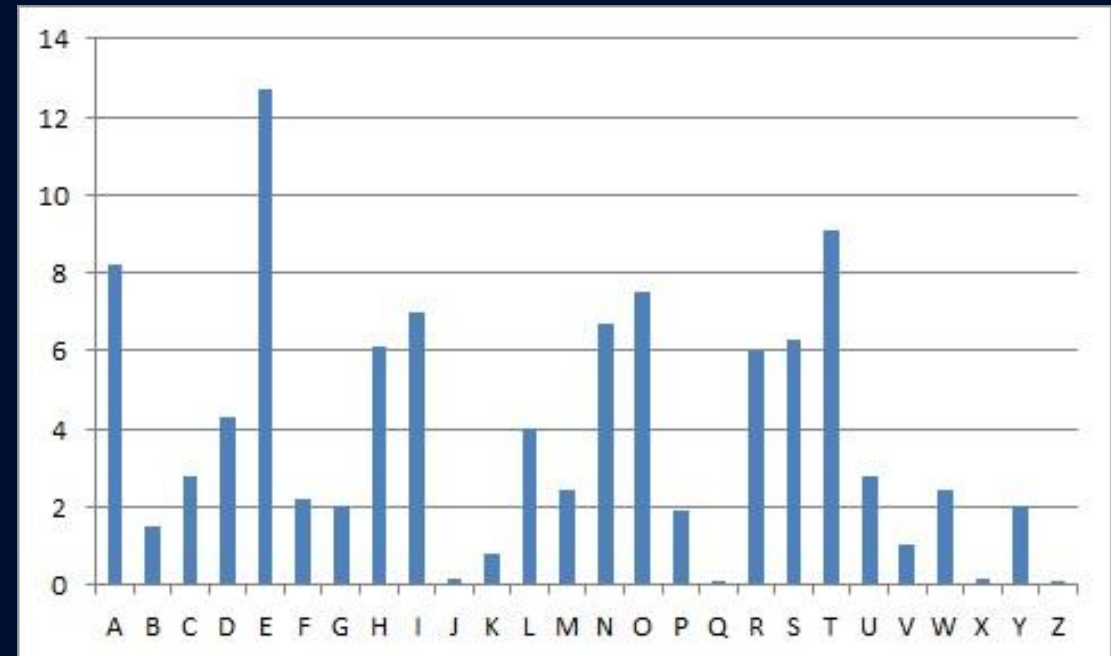
$$11 + 24 = 35 \bmod 26 = 9$$

Breaking Vigenère



- Frequency analysis
 - FACT: the most frequent letter in English: E

1. First, guess the length of the key
2. Second, count frequency
3. Third, guess the most frequent letter = E
4. Last, decrypt



Frequency analysis -> Vigenère



1. We guess the length of the key is 3

xsu nyp sre dlq cxg wiq ref bex khf
khz yvl oog xkq ref bmw kvr rvc
owm xwm xxf oxf yyq krb krb pmp
cxl skf daf orq rif khc xhc nxf oxy
vim pqy kvs pwf ovm ciy xhi swq ohr
rie bss xhz ojm bif sqq kcg xke biy
dog xkd yvy dlm ewy xhy xhm xil skf
dwg ret ofc orp ogm err sre dsw yyr
rid kfj owm pty cxy qiq krb dlc vie
orb csd kra sil dog xkq wew sqy uiq
yfm vhy cxm mvy fiy pet yyp yjw yyp
weh owr iin spm qyc dej owd bsk
dlc dlm ewy xhy xhm xil skf dw

Frequency analysis -> Vigenère



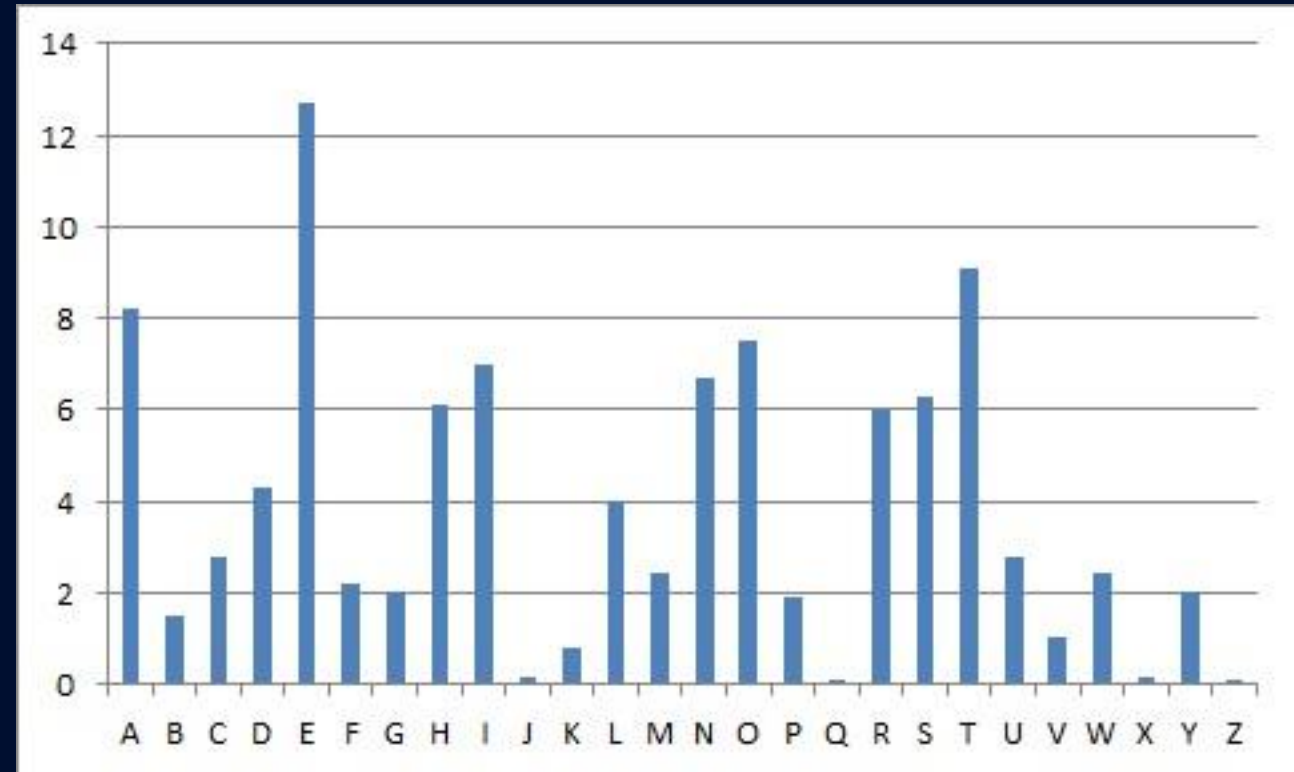
1. We guess the length of the key is 3
2. We count the letters on the 1st, 4th, 7th Spot
 1. Frequency O = 15
 2. Frequency D = 11
 3. Frequency K = 11

xsu nyp sre dlq cxg wiq ref bex khf
khz yvl oog xkq ref bmw kvr rvc
owm xwm xxf oxf yyq krb krb pmp
cxl skf daf orq rif khc xhc nxf oxy
vim pqy kvs pwf ovm ciy xhi swq ohr
rie bss xhz ojm bif sqq kcg xke biy
dog xkd yvy dlm ewy xhy xhm xil skf
dwg ret ofc orp ogm err sre dsw yyr
rid kfj owm pty cxy qiq krb dlc vie
orb csd kra sil dog xkq wew sqy uiq
yfm vhy cxm mvy fiy pet yyp yjw yyp
weh owr iin spm qyc dej owd bsk
dlc dlm ewy xhy xhm xil skf dw

Frequency analysis -> Vigenère



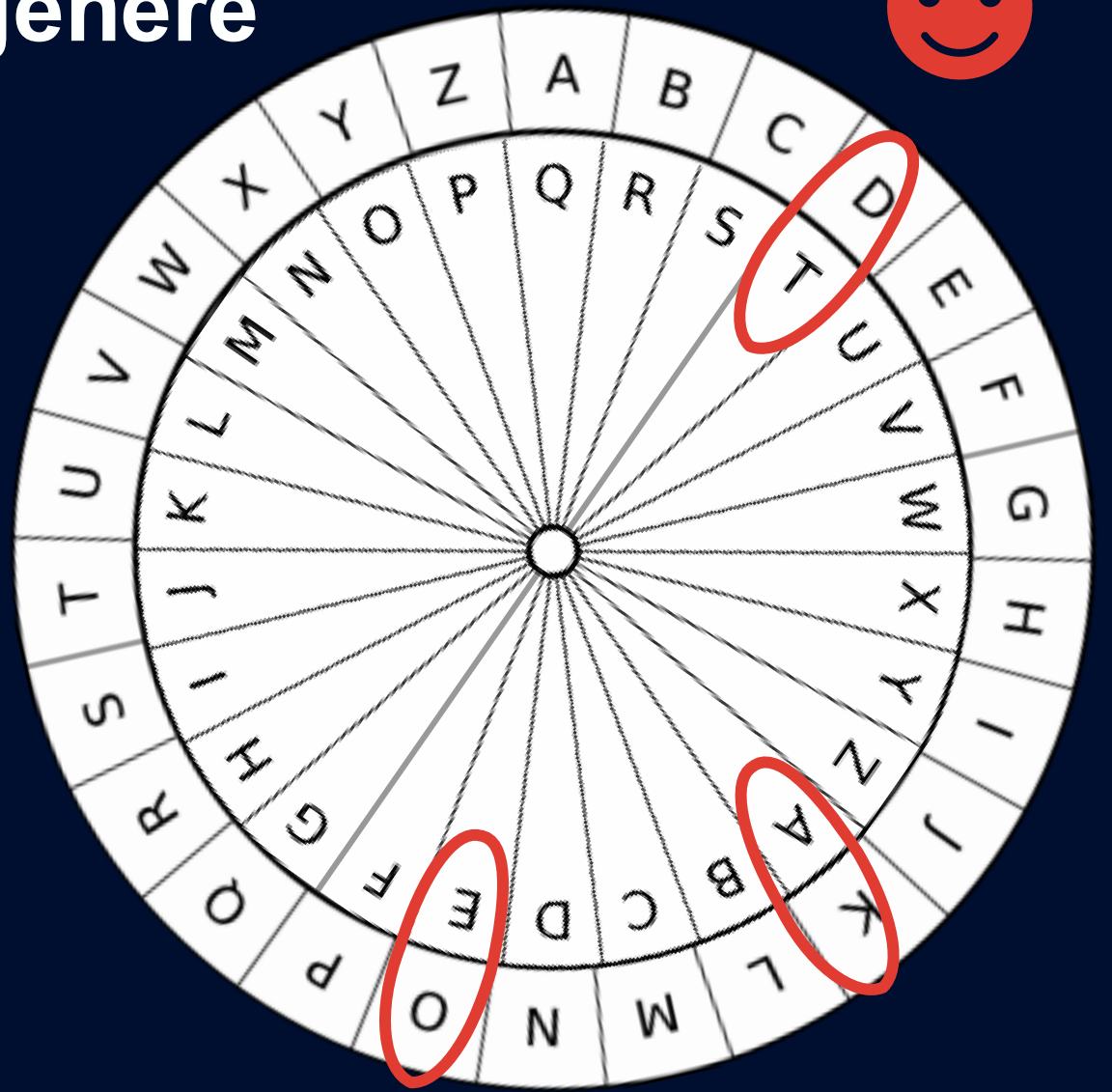
1. We guess the length of the key is 3
2. We count the letters on the 1st 4th, 7th Spot
 1. Frequency O = 15 → E
 2. Frequency D = 11 → A,T?
 3. Frequency K = 11 --> A,T?



Frequency analysis -> Vigenère



1. We guess the length of the key is 3
2. We count the letters on the 1st, 4th, 7th Spot
 1. Frequency O = 15 → E
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Frequency analysis -> Vigenère



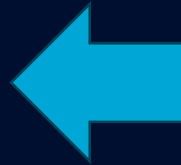
1. We guess the length of the key is **3**
2. We count the letters on the 1st 4th, 7th Spot
 1. Frequency O = 15 → **E**
 2. Frequency D = 11 → **T**
 3. Frequency K = 11 → **A**
3. Do the same for the 2nd and 3th letter of the key
4. Decrypt the message

xsu nyp sre dlq cxg wiq ref bex khf
khz yvl oog xkq ref bmw kvr rvc
owm xwm xxf oxf yyq krb krb pmp
cxl skf daf orq rif khc xhc nxf oxy
vim pqy kvs pwf ovm ciy xhi swq ohr
rie bss xhz ojm bif sqq kcg xke biy
dog xkd yvy dlm ewy xhy xhm xil skf
dwg ret ofc orp ogm err sre dsw yyr
rid kfj owm pty cxy qiq krb dlc vie
orb csd kra sil dog xkq wew sqy uiq
yfm vhy cxm mvy fiy pet yyp yjw yyp
weh owr iin spm qyc dej owd bsk
dlc dlm ewy xhy xhm xil skf dw

Frequency analysis -> Vigenère



Now during this time Shahrazad had borne king Shahriyar three sons. On the thousand and first night, when she had ended the tale of Ma'aruf, she rose and kissed the ground before him, saying: “great king, for a thousand and one nights I have been recounting to you the fables of past ages and the legends of ancient kings. May I make so bold as to crave a favour of your majesty?”
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xsu nyp sre dlq cxg wiq ref bex khf
khz yvl oog xkq ref bmw kvr rvc
owm xwm xxf oxf yyq krb krb pmp
cxl skf daf orq rif khc xhc nxf oxy
vim pqy kvs pwf ovm ciy xhi swq ohr
rie bss xhz ojm bif sqq kcg xke biy
dog xkd yvy dlm ewy xhy xhm xil skf
dwg ret ofc orp ogm err sre dsw yyr
rid kfj owm pty cxy qiq krb dlc vie
orb csd kra sil dog xkq wew sqy uiq
yfm vhy cxm mvy fiy pet yyp yjw yyp
weh owr iin spm qyc dej owd bsk
dlc dlm ewy xhy xhm xil skf dw



Real world example

Enigma | World war I and II

Explanation video Enigma:

<https://www.youtube.com/watch?v=ybkkiGtJmkM>

Movie of Alan Turing:

The Imitation Game

Invention of Enigma



- In 1915
- By two Dutch for the Dutch War Ministry
- First prototype by the end of WW I

What is enigma



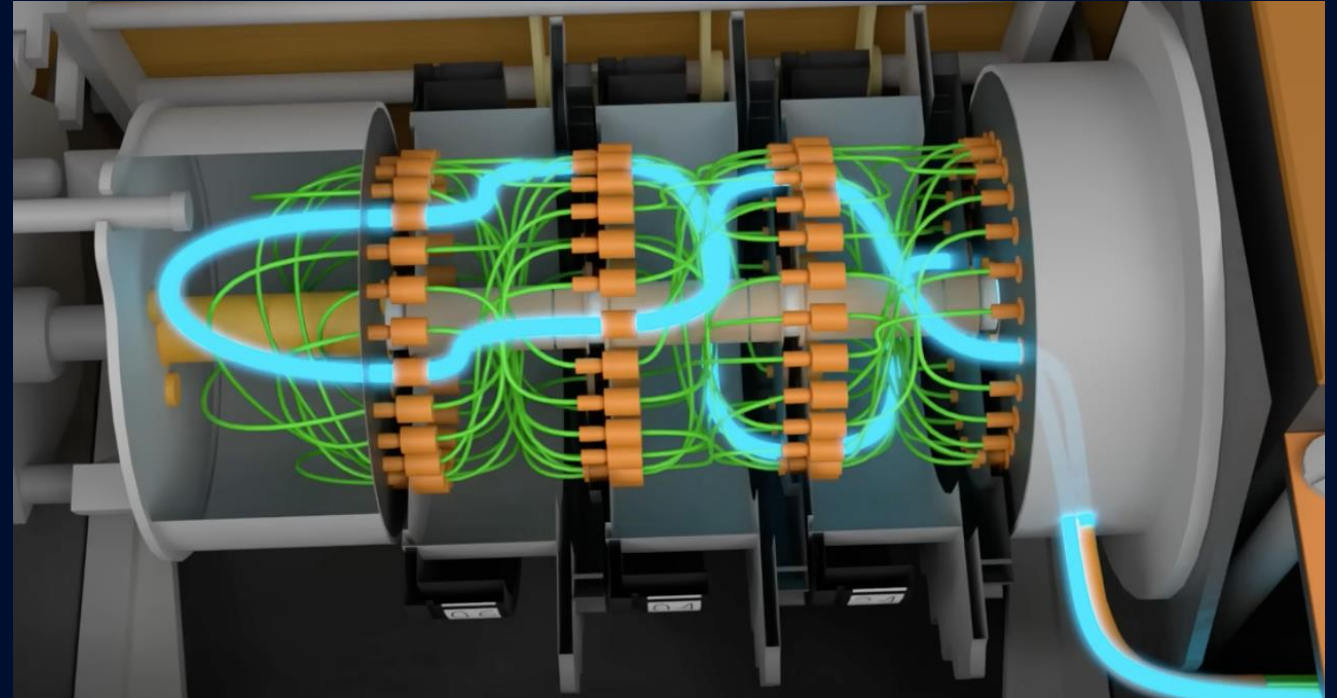
- Machine to encrypt/ decrypt
- Keyboard to type text
- Lightbulbs to see encrypted text
- Encryption mechanism
 - 3 wheels → encrypt letters



What is enigma



- Machine to encrypt/ decrypt
- Keyboard to type text
- Lightbulbs to see encrypted text
- Encryption mechanism
 - 3 wheels → encrypt letters
 - A lot of wires



What is enigma



- Machine to encrypt/ decrypt
- Keyboard to type text
- Lightbulbs to see encrypted text
- Encryption mechanism
 - 3 wheels → encrypt letters
 - A lot of wires
 - “Stecker” → swap letters



Encryption – Decryption with Enigma



- Encryption and decryption was done with the same machine
- But both machines need the same setting:
 - The orientation of the wheels
 - The “steckers” in the stecker board
- Distribute these settings via:
 - Sheets
 - Codebooks



Evolution of Enigma

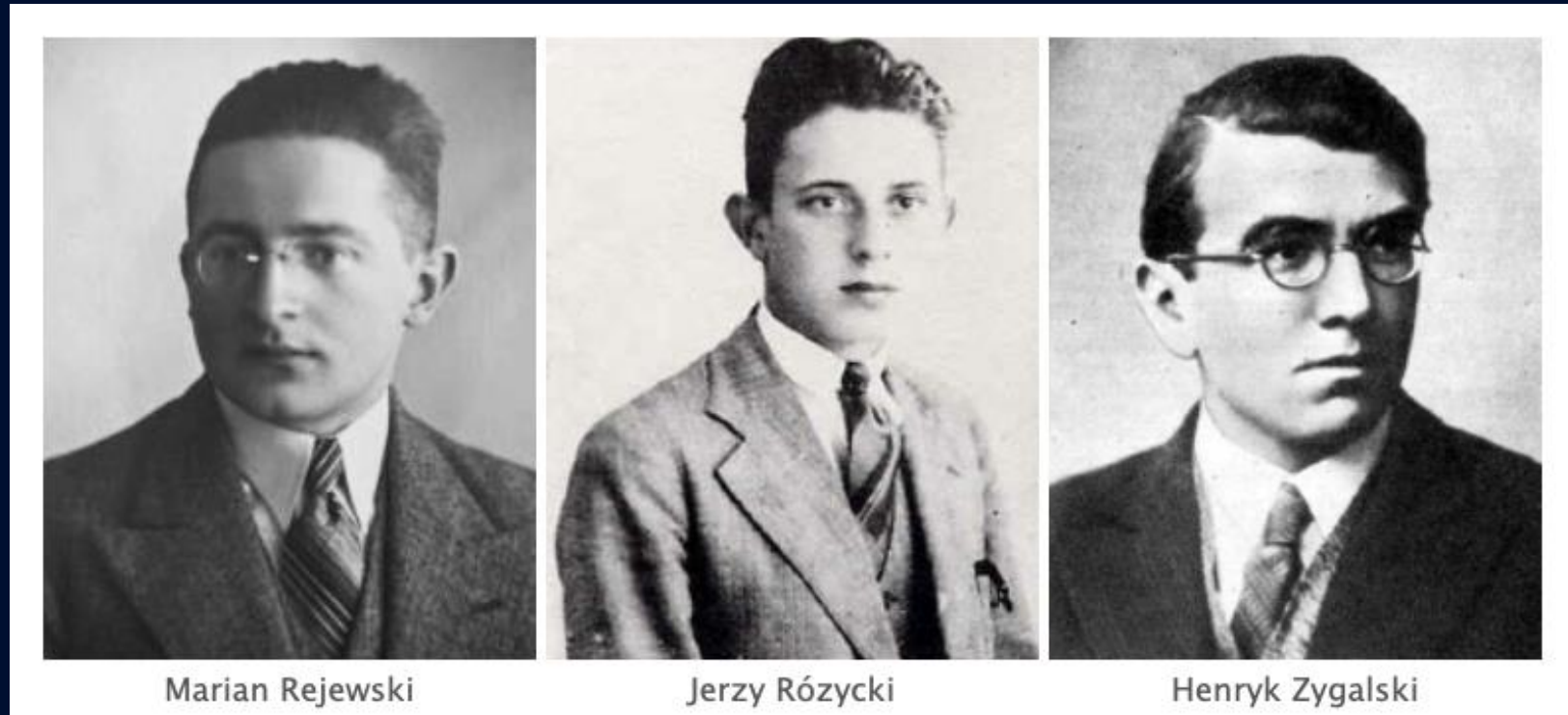
- Enigma A (1924) → 26 keys
- Enigma B (1924) → 28 keys
- Enigma C (1925) → 29 keys
- Enigma D (1926)

- Enigma K (1927) → Increased cipher strength
- Enigma I (1927-1929) → Exclusively used for military purposes

Military Enigma (Enigma I) Broken



- 3 Polish mathematicians broke Enigma I
- They needed around 100 intercepted messages to recover the initial settings



Marian Rejewski

Jerzy Rózycki

Henryk Zygalski

Stronger Enigma

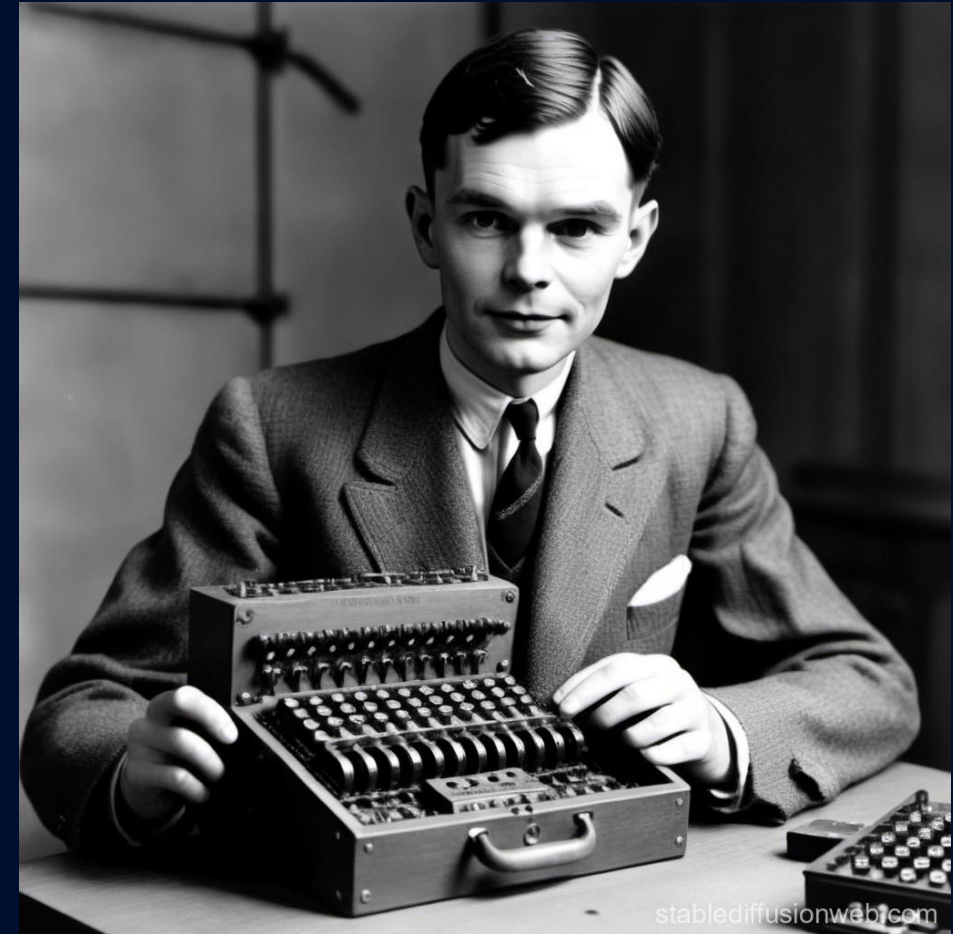


- Every enigma has an initial setting
 - 3 wheels
 - Stecker board
- Germans used only this setting
- But the Polish broke it with 100 messages... so...
- On **15 September 1938** the Germans had a new setting every day
- 2 wheels were added
 - 10 times more possible settings

Bletchley Park



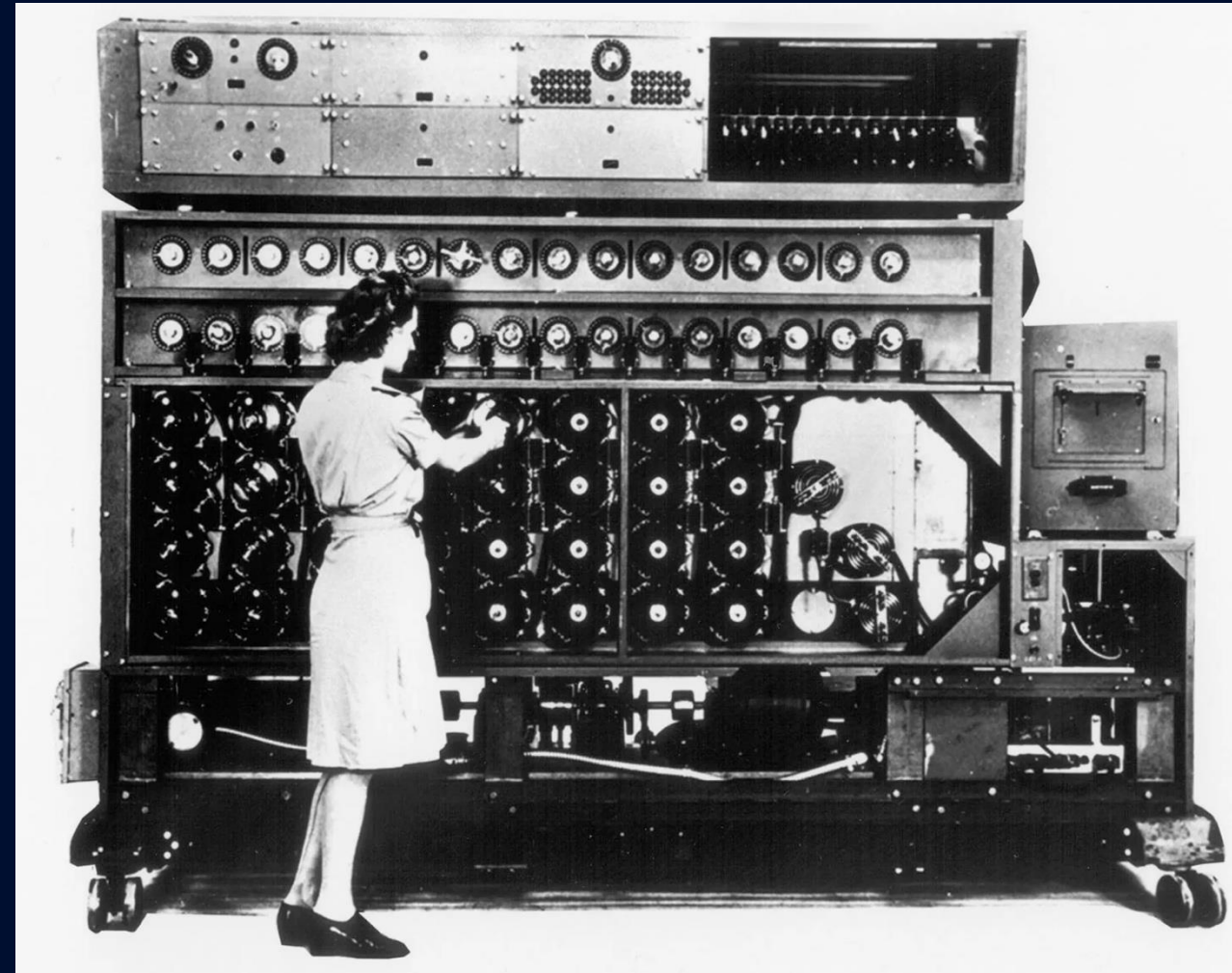
- 45 miles north of London
- Since **14 August 1939** hosted Government Code and Cypher School
- Staff:
 - Professional code breakers
 - Chess players
 - Mathematicians
 - Organizers
- Among them: **Alan Turing**



Breaking Enigma



- Messages were broken by hand
- Alan Turing made the Bombe; a machine to break messages



New Enigma



- On **1 February 1942**, the German Navy introduced a new Enigma machine
- New wheel
- New indicator system
- New codebooks



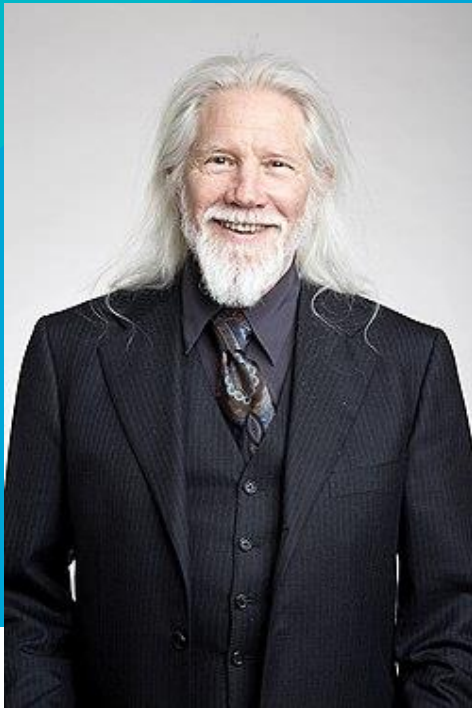
Breaking New Enigma

- Alan Turing discovered the new wiring
- They captured a new codebook **30 October 1942**
- But Bombe was not suitable for breaking these machines
- The UK did not have enough resources to build a new machine
- The US had sufficient supplies and build the US Bombe

Why the Germans kept using Enigma



- Germany always believed that Enigma could not be broken
- The German Navy, however, had their doubt
 - The Navy thus introduced extra security measurements
- The German Air Force also added security
 - An extra device to be attached to Enigma
 - However, “due to operator mistakes it was broken within a few days after its introduction”
- Even after the war, it was kept secret that Enigma was broken







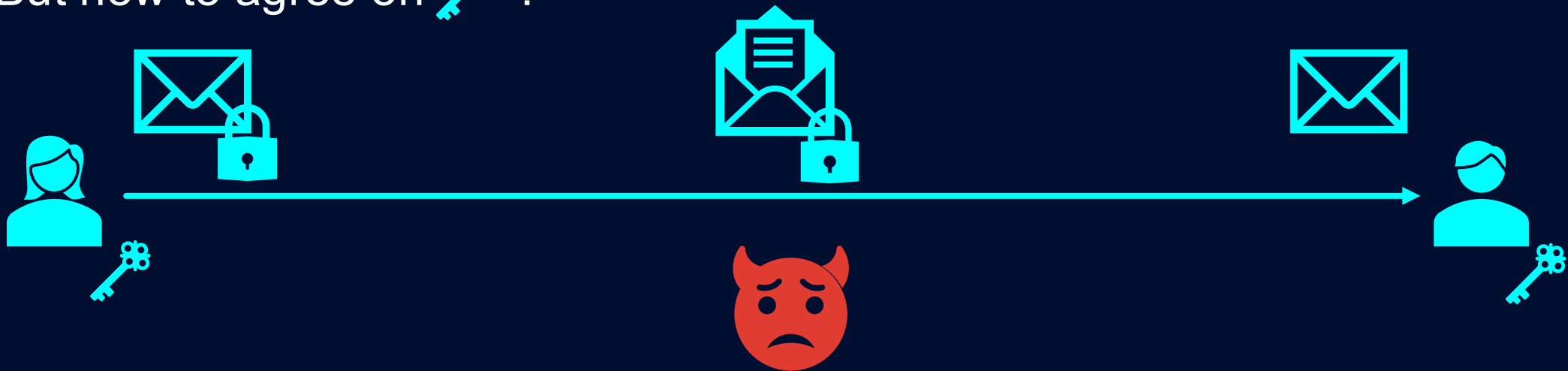
Diffie-Hellman

1976 – Building block of current cryptography

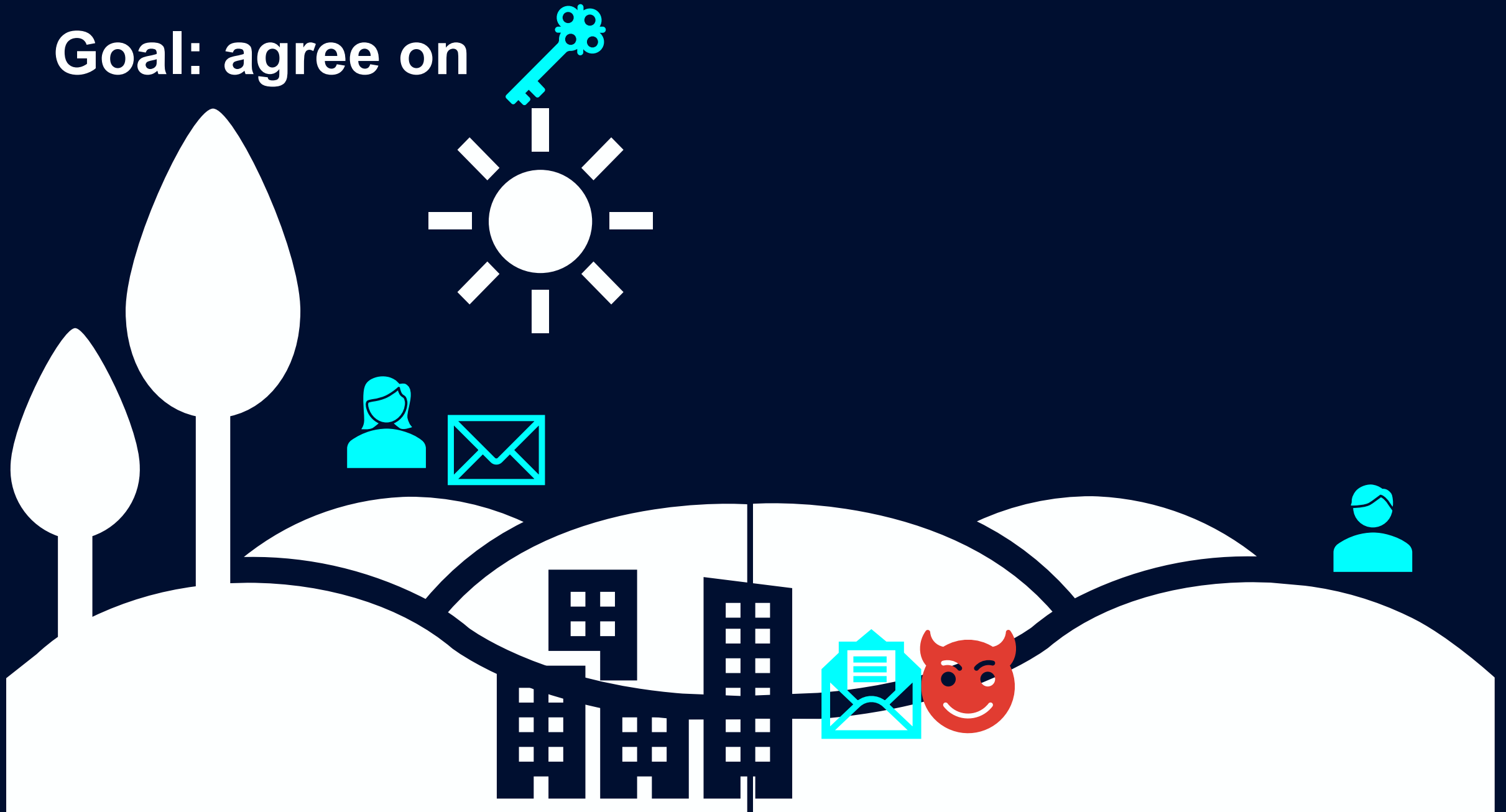


Recall: Goal of Alice

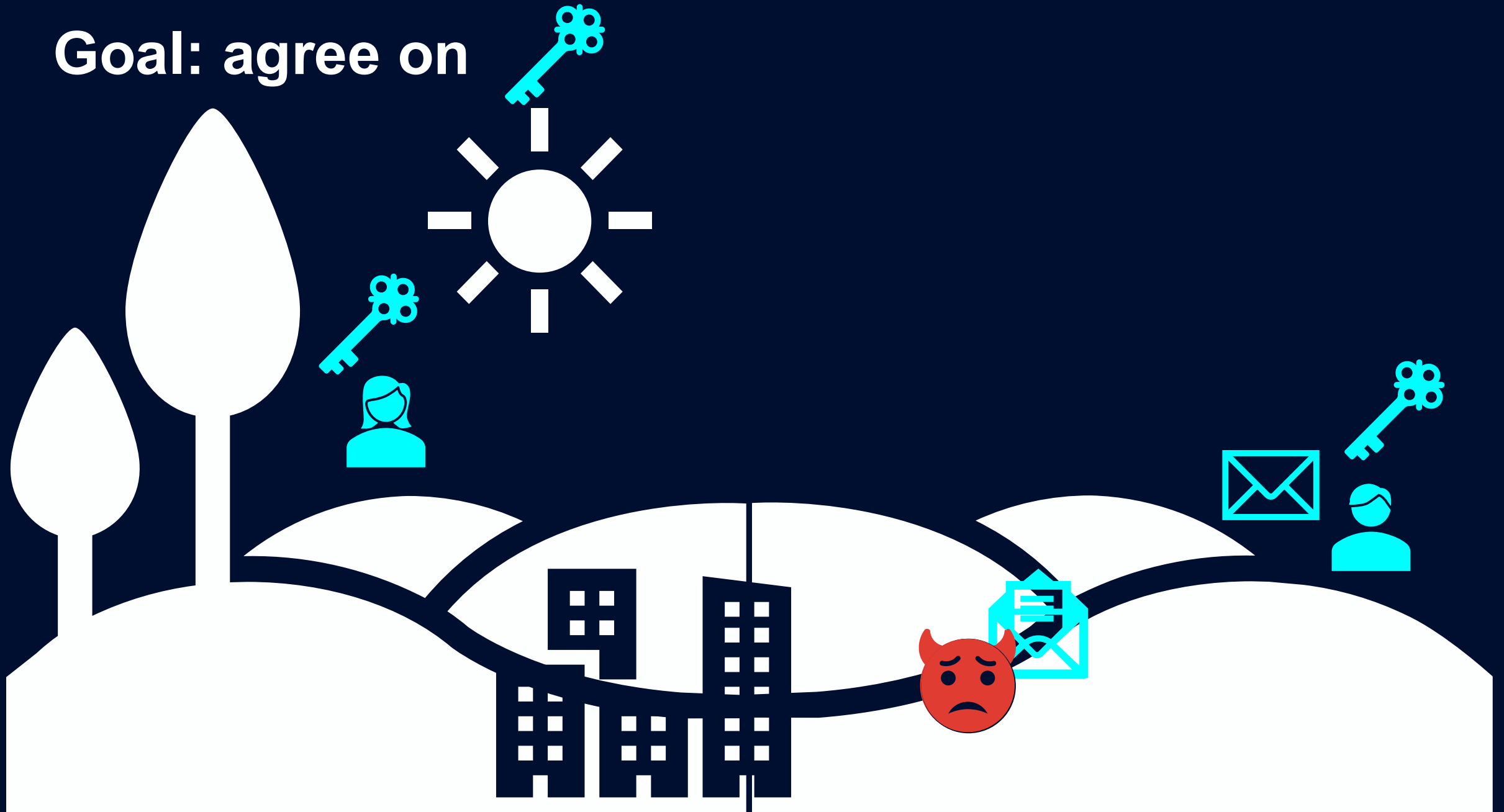
- A key  that only Alice and Bob know
- Alice can *Encrypt* the message 
- Bob can *Decrypt* the message 
- If Eve sees the message, she will not understand
- But how to agree on  ?



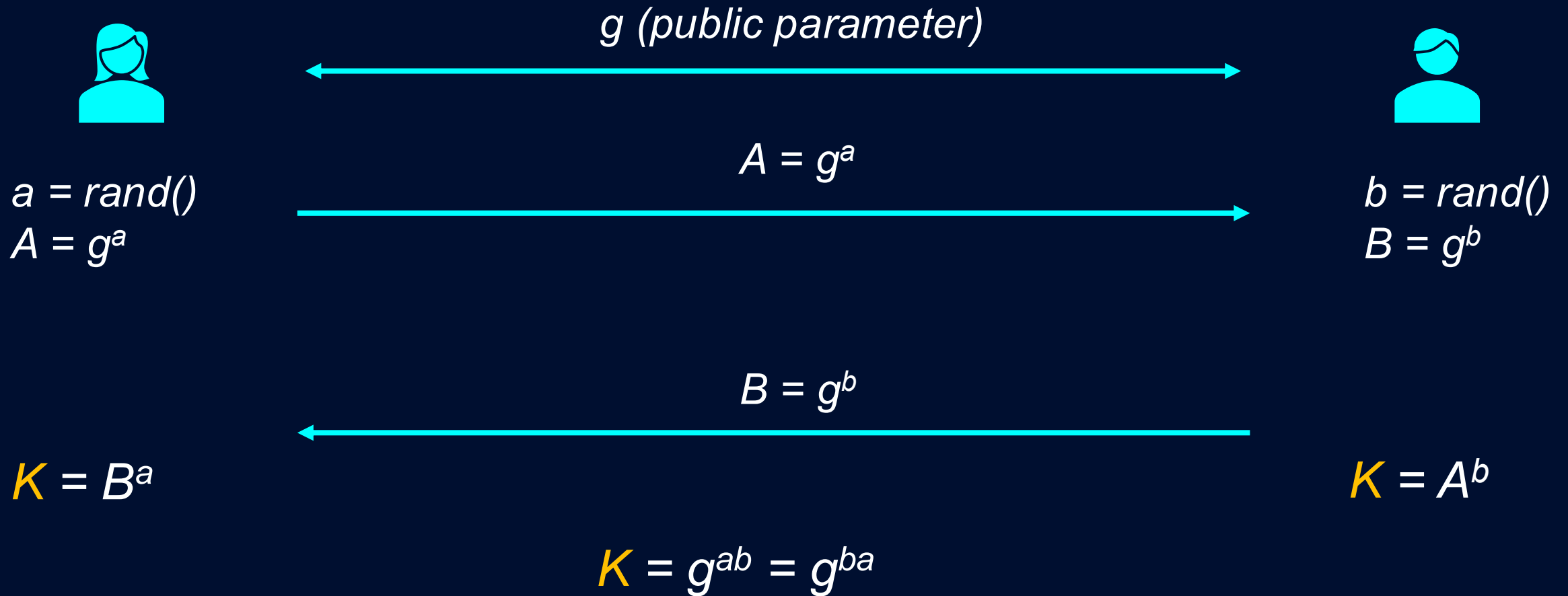
Goal: agree on



Goal: agree on



Diffie Hellman exchange

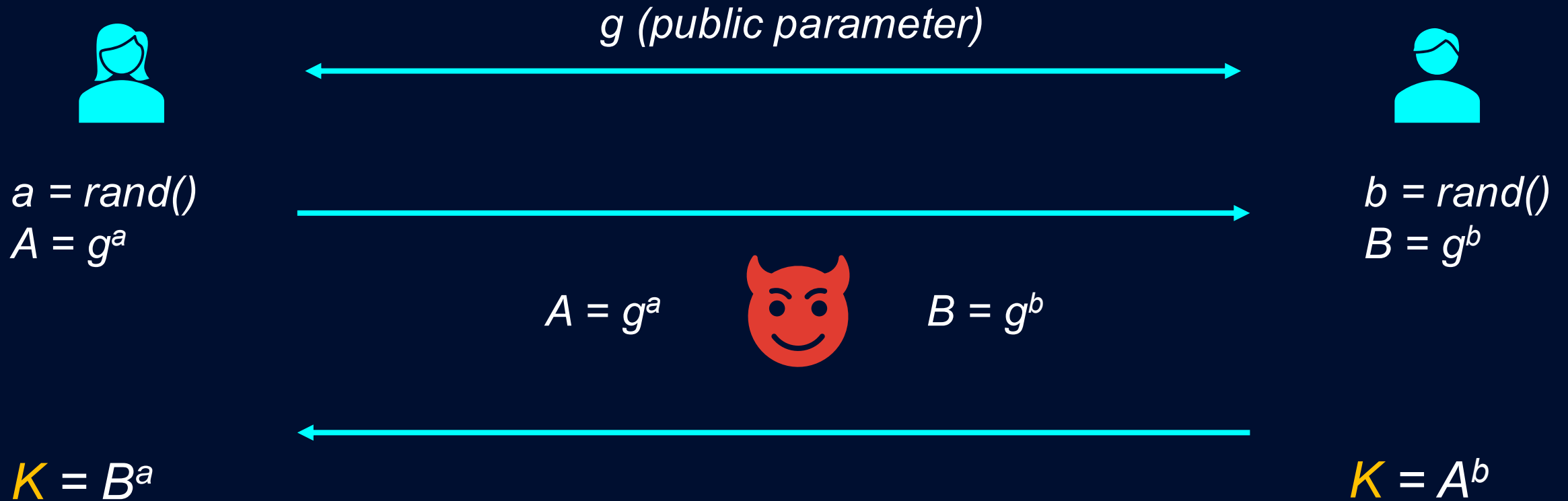


Diffie Hellman



exchange

INSECURE



$$a = \log_g(A)$$
$$K = B^a$$

“Normal” arithmetic

- Given a, g
- **Easy** to compute:

$$A \equiv g^a$$

- Given A, g
- **Easy** to compute:

$$a \equiv \log_g(A)$$

Modular arithmetic

- Given a, g, p
- **Easy** to compute:

$$A \equiv g^a \bmod p$$

- Given A, g, p
- **Difficult** to compute:

$$a \equiv \log_g(A) \bmod p$$

Diffie Hellman



exchange

INSECURE



g (public parameter)



$$a = \text{rand}()$$
$$A = g^a$$

$$b = \text{rand}()$$
$$B = g^b$$



$$K = B^a$$

$$K = A^b$$

$$K = g^{ab} = g^{ba}$$

Diffie Hellman



exchange

SECURE



g and p (public parameter)



$$a = \text{rand}()$$
$$A = g^a \text{ mod } p$$

$$b = \text{rand}()$$
$$B = g^b \text{ mod } p$$



$$\text{key} = B^a \text{ mod } p$$

$$\text{key} = A^b \text{ mod } p$$

$$K = g^{ab} = g^{ba}$$

Diffie Hellman



exchange

SECURE



g and p (public parameter)



$$a = \text{rand}()$$
$$A = g^a \text{ mod } p$$

$$b = \text{rand}()$$
$$B = g^b \text{ mod } p$$

$$A = g^a$$



$$B = g^b$$

$$\text{key} = B^a \text{ mod } p$$

$$\text{key} = A^b \text{ mod } p$$

~~$$a = \log_g(A)$$~~

Breaking Diffie-Hellman





Thanks for your
attention!